



**WINTECHCON**

## **Kernelization of vertex cover based on crown structure**

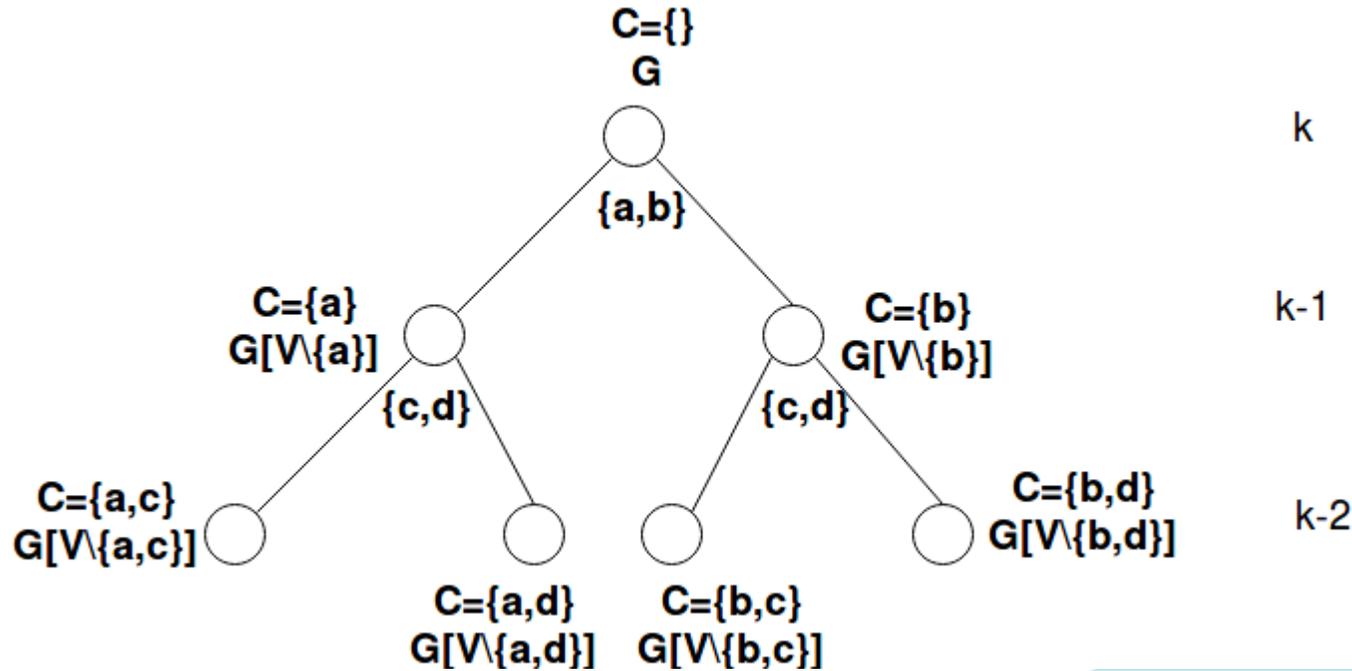
September 27, 2019

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# Vertex Cover problem

Given a graph  $G=(V,E)$ , find a subset  $C$  of  $V$  of size at most  $k$  such that for each edge  $\{u,v\}$  in  $E$ , either  $u$  or  $v$  belongs to  $C$



Time complexity of this solution is  $O((n+m)2^k)$ .

# Fixed parameter tractable algorithms

- Instead of expressing the running time as a function  $T(n)$  of  $n$ , we express it as a function  $T(n,k)$  of the input size  $n$  and some parameter  $k$  of the input
- Algorithms whose running times are exponential only in a chosen parameter, and are polynomial in the instance size
- $T(n,k) = O(\text{poly}(n)f(k))$  where  $f()$  is an exponential function and  $\text{poly}()$  is a polynomial function
- Efficient on all inputs where  $k$  is small.



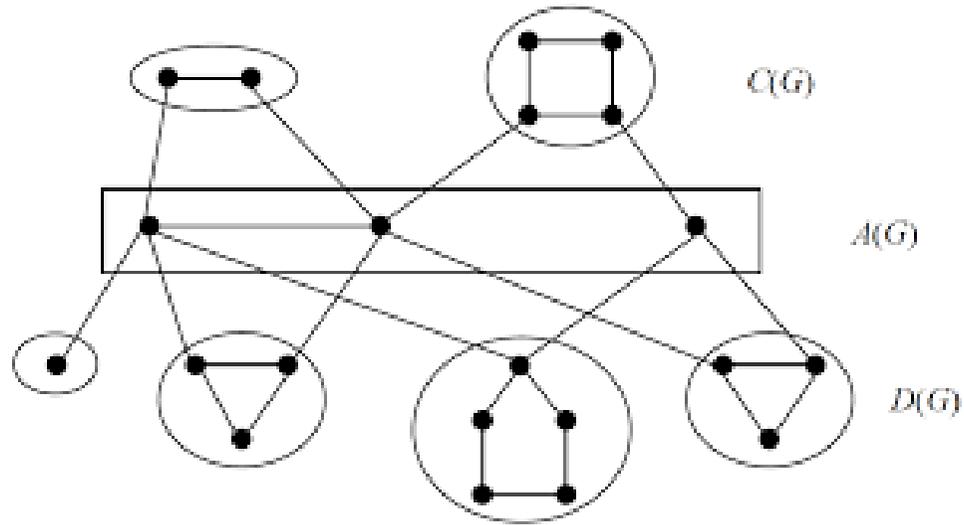
# Kernelization

- Reduce problem instance  $(I, k)$  to  $(I', k')$
- Size of  $I'$  is  $\text{poly}(k)$
- Time to solve reduced instance is  $\text{poly}(|I'|, k')$
- $k' \leq k$
- Reduction takes polynomial time
- Reduced instance has a solution iff given instance has a solution



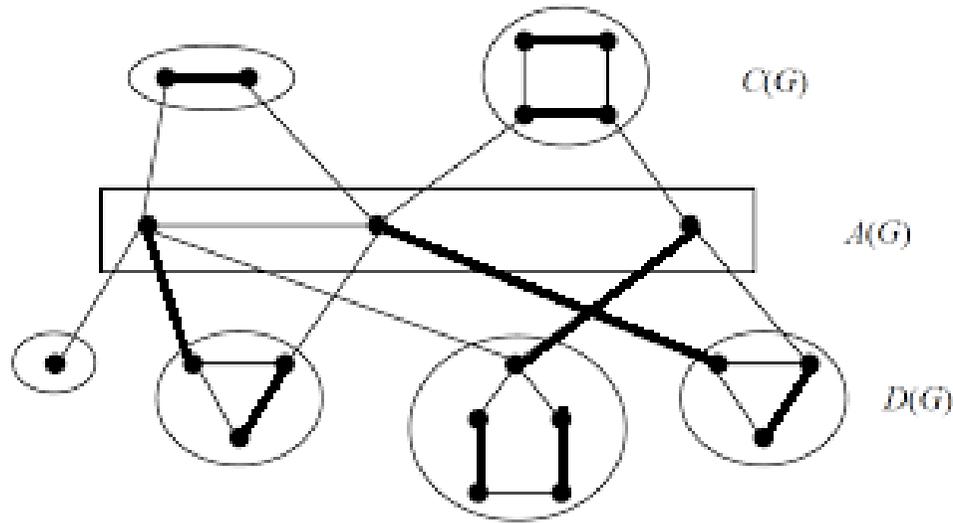
# Kernelization method

- The goal is to reduce the problem instance for vertex cover problem
- The reduced instance will have at most  $2k$  vertices
- Gallai-Edmonds decomposition



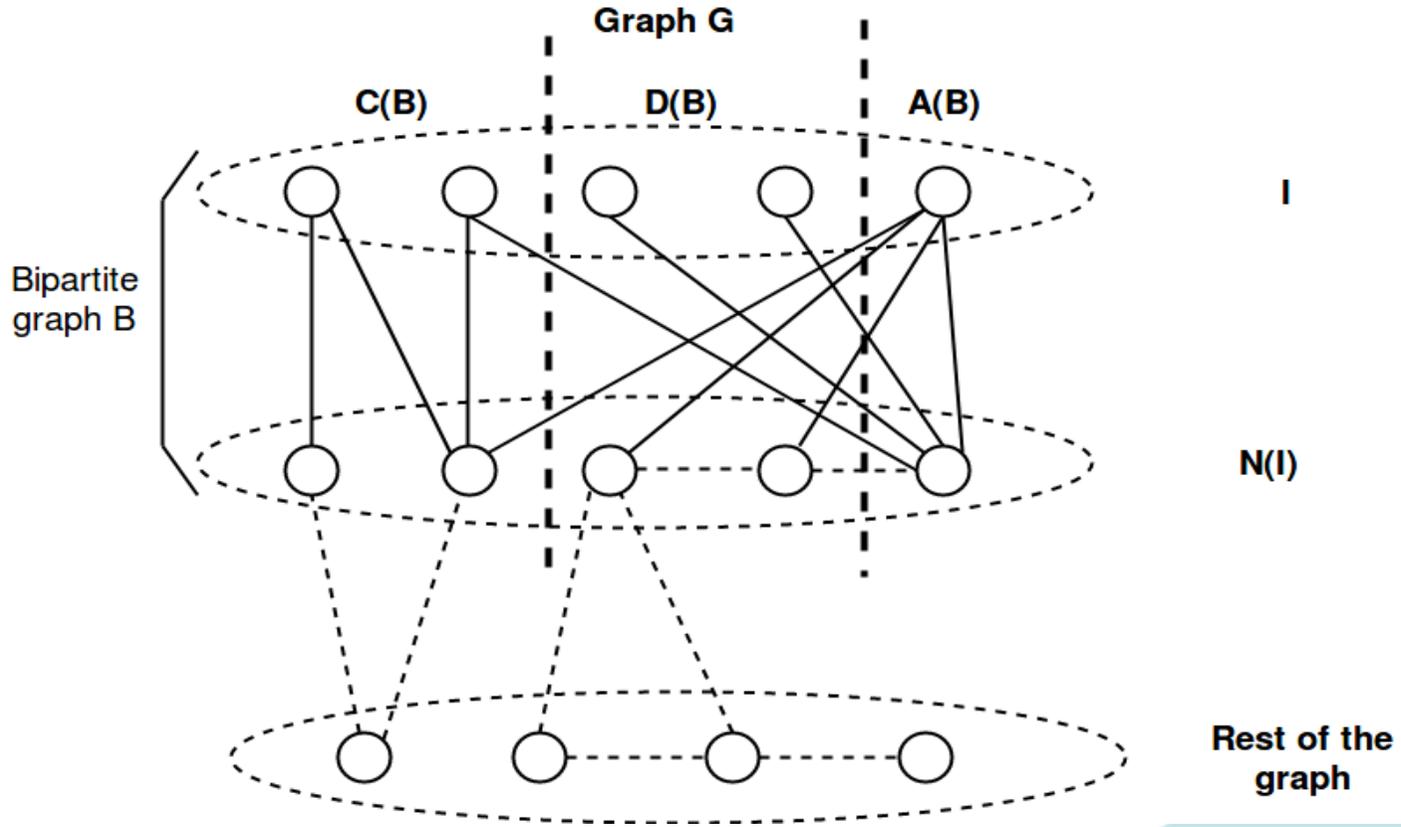
Source: Graphs of Edge-Intersecting and Non-Splitting Paths PART II: One-Bend ENPG

# Properties of Gallai-Edmonds decomposition



- $C(G)$  and  $A(G)$  vertices are matched in all maximum matchings
- $A(G)$  vertices are matched to  $D(G)$  vertices and  $C(G)$  is perfectly matched in a maximum matching
- $D(G)$  components are near-perfectly matchable. One vertex is unmatched or matched outside.
- $D(G)$  is an independent set in bipartite graphs

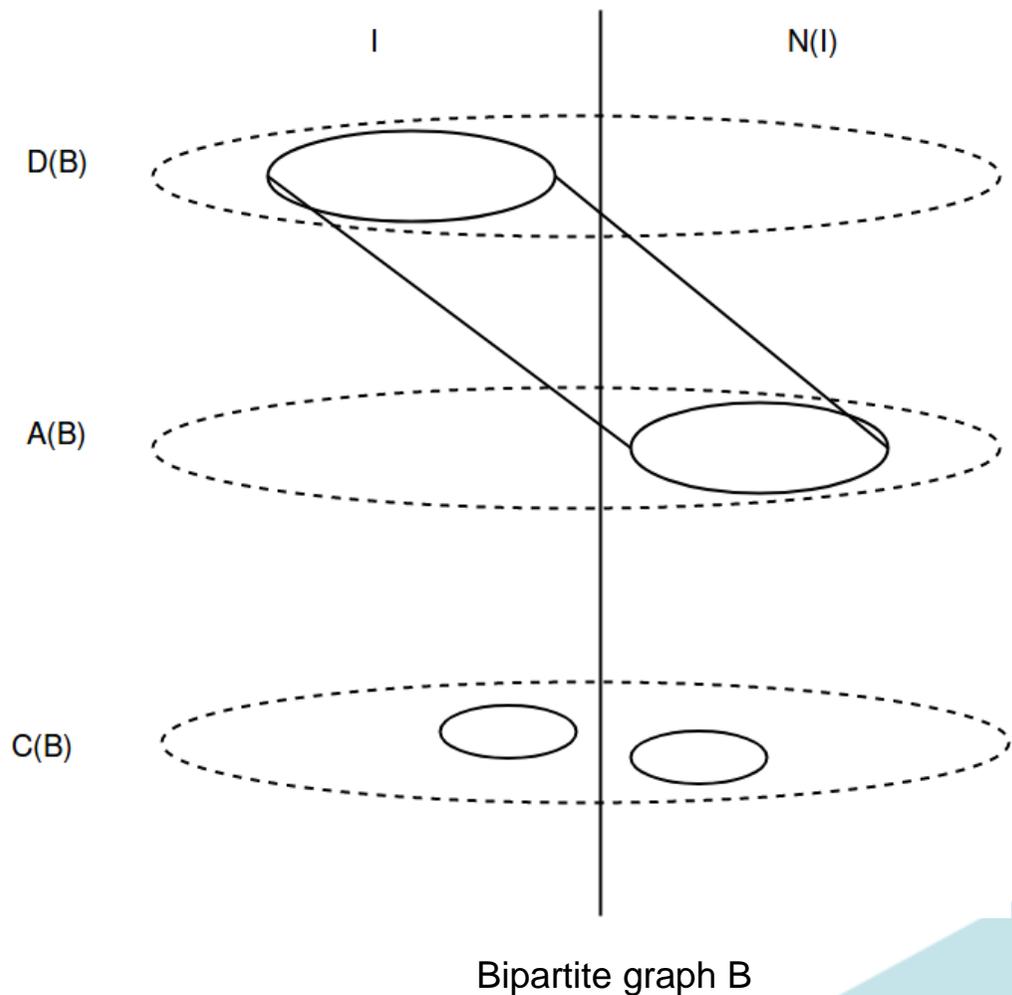
# Reduction



Set of isolated vertices of  $D(G)$  and their neighbors denoted by  $I$  and  $N(I)$  respectively



# Reduction

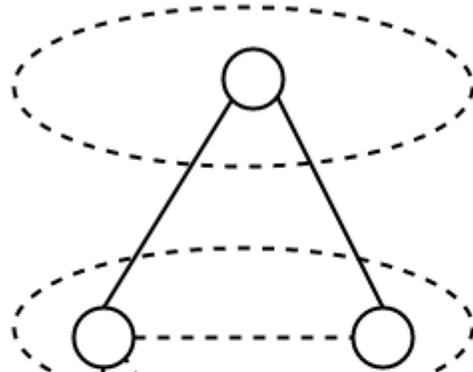


There exists a minimum vertex cover of  $G$  which does not contain any vertex from  $(C(B) \cup D(B)) \cap I$  and contains all vertices from  $(C(B) \cup A(B)) \cap N(I)$ .



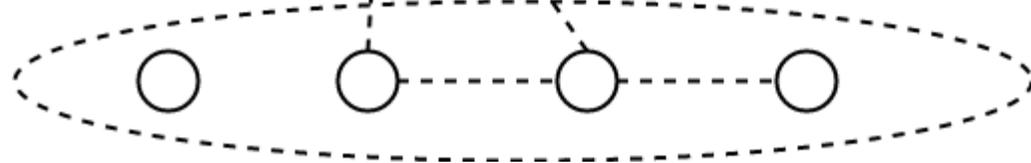
# Reduction

Graph  $G'$



$$I' = A(B)$$

$$N(I') = D(B)$$



Reduced Graph  $G'$

Rest of  
 $G'$

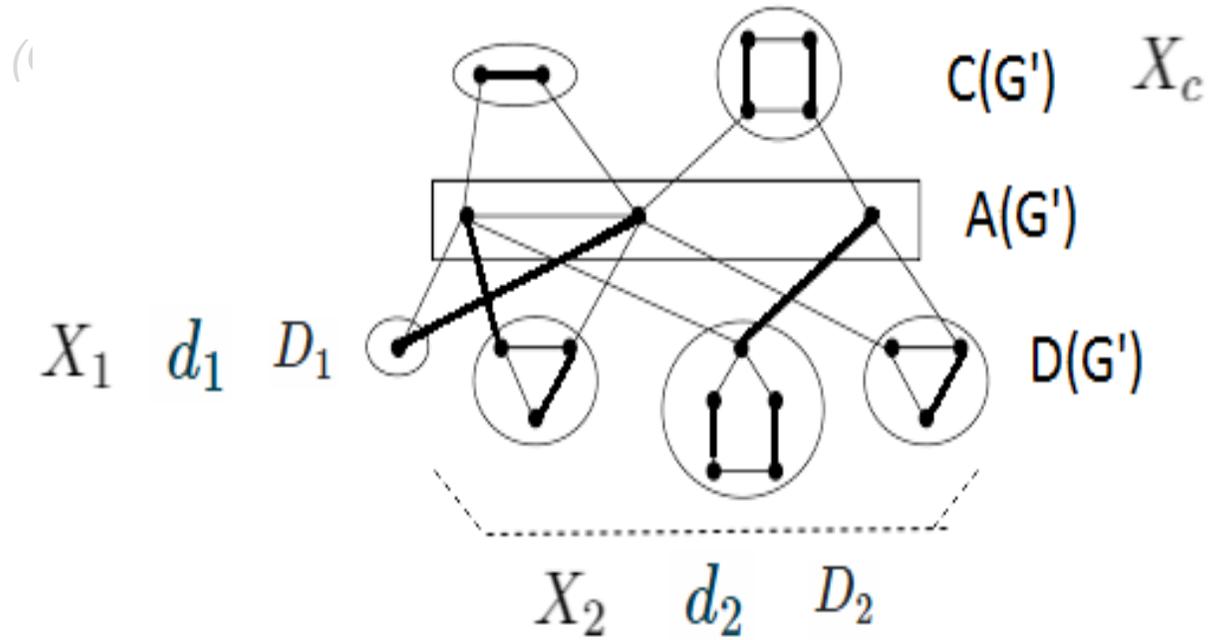
*experimental work)*

If  $C'$  is a minimum vertex cover of the reduced graph  $G'$ , then  $C = C' \cup [(C(B) \cup A(B)) \cap N(I)]$  is a minimum vertex cover of  $G$ .

There exists a maximum matching for  $G'$  in which all  $I'$  vertices are matched.



# 2k kernel

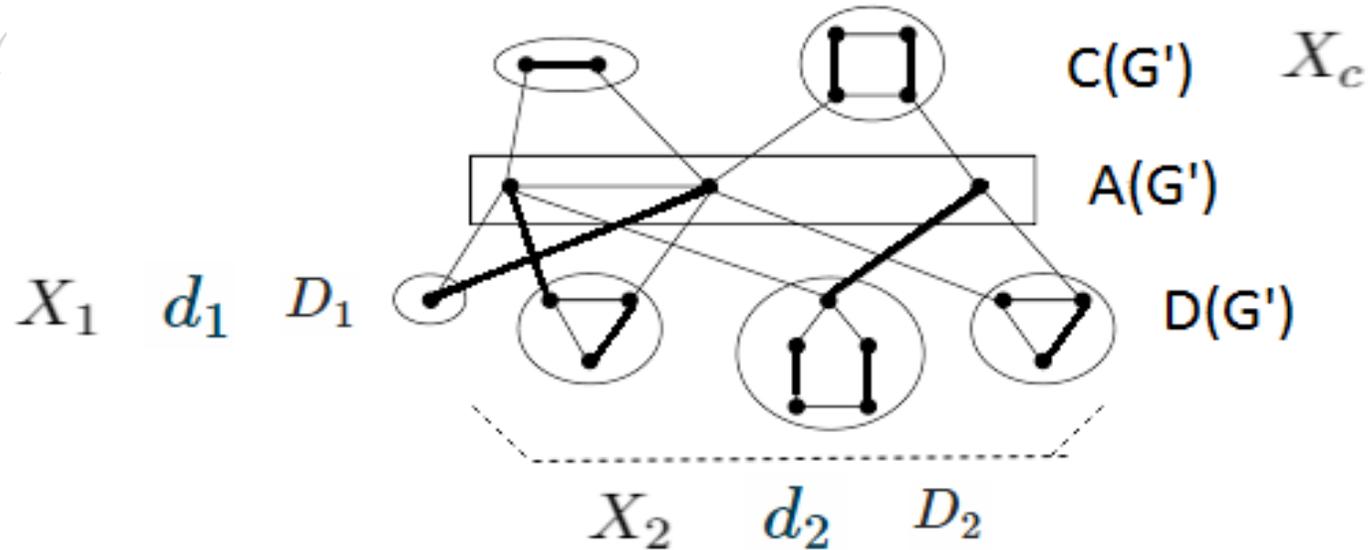


Gallai-Edmonds decomposition of  $G'$

Let  $M$  be the maximum matching for  $G'$  in which all vertices of  $I'$  are matched.



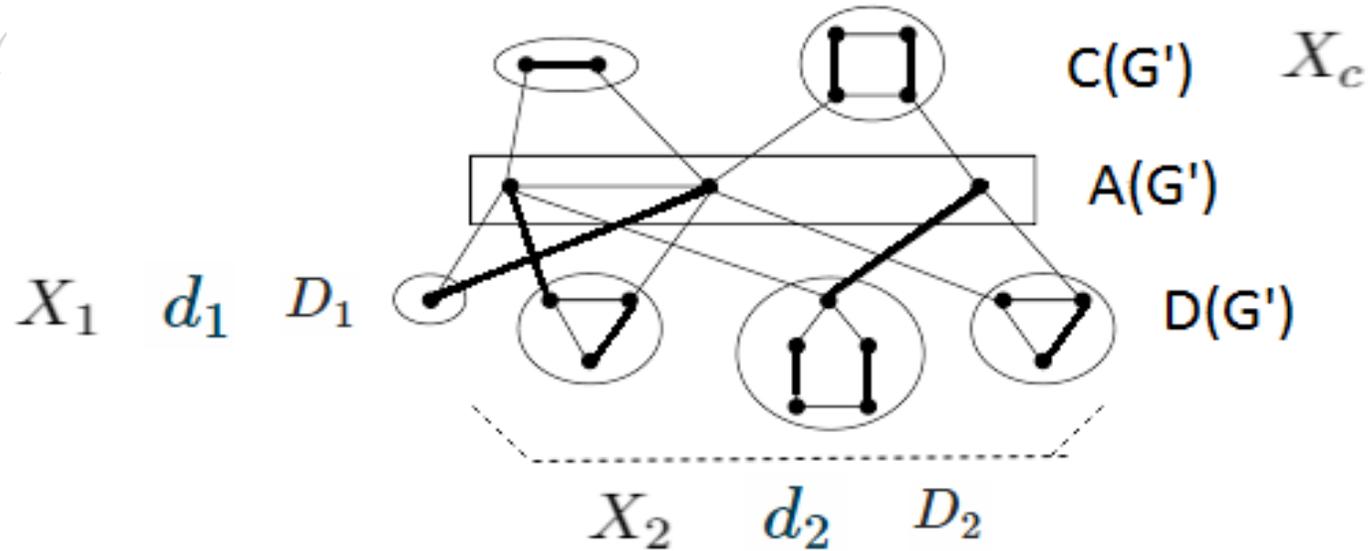
## 2k kernel



Gallai-Edmonds decomposition of  $G'$

- $|A(G')| \leq D_1 + D_2$  from properties of Gallai-Edmonds decomposition
- $|A(G')| + |D(G')| \leq D_1 + D_2 + d_1 + d_2$
- $X_1 \geq d_1$  because in  $M$ , all  $d_1$  vertices are matched
- $X_2 \geq (d_2 + D_2)/2$  because for any component of  $D(G')$  with size  $2r+1$  where  $r \geq 1$ , at least  $r+1$  vertices are required to cover the vertex cover

## 2k kernel



Gallai-Edmonds decomposition of  $G'$

- $|A(G')| + |D(G')| \leq D_1 + D_2 + d_1 + d_2 = 2D_1 + 2(D_2 + d_2)/2 \leq 2(X_1 + X_2)$
- $|C(G')| \leq 2|X_c|$  because number of matching edges of  $C(G')$  is exactly  $|C(G')|/2$
- $|V| = |D(G')| + |A(G')| + |C(G')| \leq 2(|X_1| + |X_2| + |X_c|) \leq 2|X|$

# Conclusions

*(Click to add text)*

- An algorithm was presented for reducing the problem instance for vertex cover to a graph of size  $2k$ , where  $k$  is size of vertex cover
- Finding a minimum vertex cover find its applications in keyword based text summarization.
- Simulating propagation of computer worms on a network and designing techniques to prevent them also make use of minimum vertex cover



# References

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